

A Trust Funnel Algorithm for Nonconvex Equality Constrained Optimization with $\mathcal{O}(\epsilon^{-3/2})$ Complexity

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OPTIMIZATION AND ITS APPLICATIONS

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Outline

Motivation

Proposed Algorithm

Theoretical Results

Numerical Results

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Introduction

Consider nonconvex equality constrained optimization problems of the form

$$\begin{aligned} \min_{x \in \mathbb{R}^n} f(x) \\ \text{s.t. } c(x) = 0. \end{aligned}$$

where $f : \mathbb{R}^n \rightarrow \mathbb{R}$ and $c : \mathbb{R}^n \rightarrow \mathbb{R}^m$ are twice continuously differentiable.

- ▶ We are interested in algorithm worst-case iteration / evaluation complexity.
- ▶ Constraints are not necessarily linear!

Algorithms for equality constrained (nonconvex) optimization

Sequential Quadratic Programming (SQP) / Newton's method

Trust Funnel; Gould & Toint (2010)

Short-Step ARC; Cartis, Gould, & Toint (2013)

Algorithms for equality constrained (nonconvex) optimization

Sequential Quadratic Programming (SQP) / Newton's method

- ▶ **Global convergence:** globally convergent (trust region/line search)

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Algorithms for equality constrained (nonconvex) optimization

Sequential Quadratic Programming (SQP) / Newton's method

- ▶ **Global convergence**: globally convergent (trust region/line search)
- ▶ **Worst-case complexity**: No proved bound

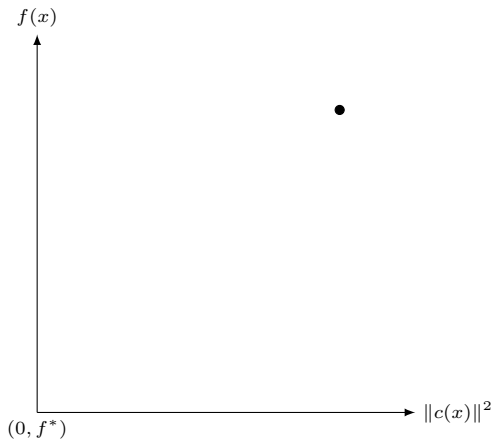
Trust Funnel; Gould & Toint (2010)

- ▶ **Global convergence**: globally convergent
- ▶ **Worst-case complexity**: No proved bound

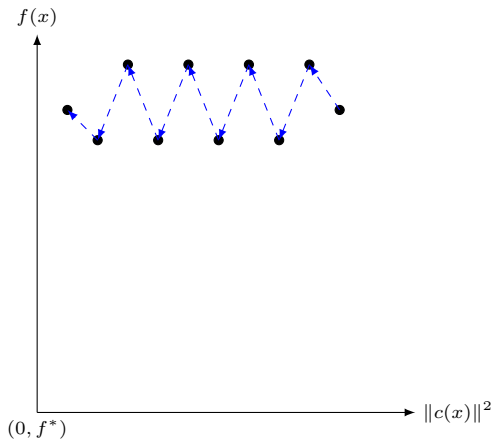
Short-Step ARC; Cartis, Gould, & Toint (2013)

- ▶ **Global convergence**: globally convergent
- ▶ **Worst-case complexity**: $\mathcal{O}(\epsilon^{-3/2})$

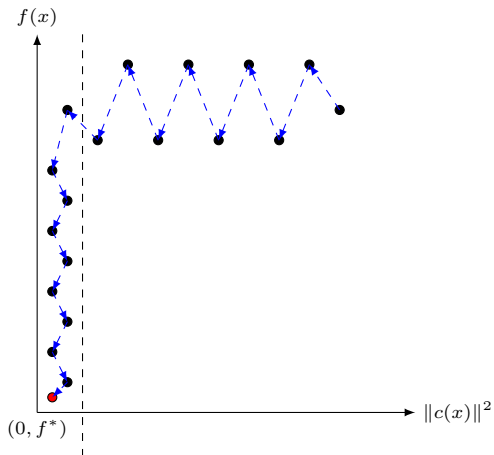
Short-Step ARC



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Main Concerns

- ▶ Completely ignores the objective function during the first phase
- ▶ **Question:** Can we do better?

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- ▶ Completely ignores the objective function during the first phase
- ▶ **Question:** Can we do better?
- ▶ **Yes!(?)**
- ▶ First, rather than two-phase approach that ignores objective in phase 1, wrap in a **trust funnel** framework that observes objective in both phases.
- ▶ Second, consider TRACE method for unconstrained nonconvex optimization
 - ▶ F. E. Curtis, D. P. Robinson, MS, “A trust region algorithm with a worst-case iteration complexity of $\mathcal{O}(\epsilon^{-3/2})$ for nonconvex optimization,” *Mathematical Programming*, 162, 2017.

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SQP “core”

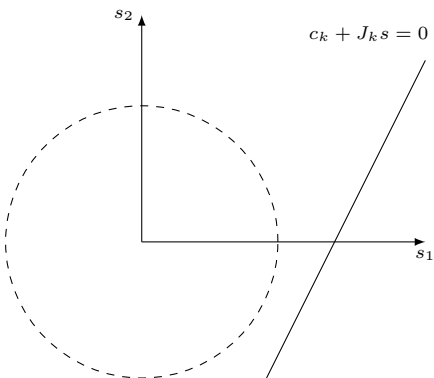
- ▶ Given x_k , find s_k as a solution of

$$\begin{aligned} \min_{s \in \mathbb{R}^n} \quad & f_k + g_k^T s + \frac{1}{2} s^T H_k s \\ \text{s.t.} \quad & c_k + J_k s = 0 \end{aligned}$$

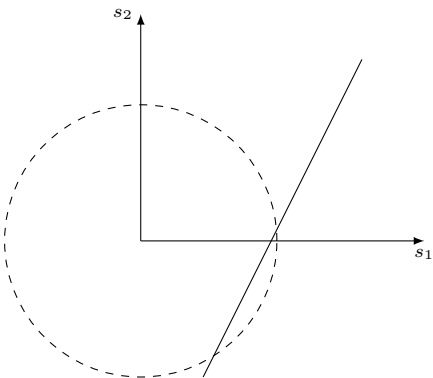
Issues:

- ▶ H_k might not be positive definite over $\text{Null}(J_k)$.
- ▶ Trust region! . . . but constraints might be incompatible.

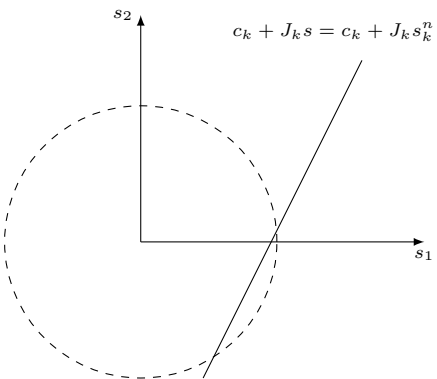
Step decomposition



Step decomposition



Step decomposition



Trust funnel basics

Step decomposition approach:

- ▶ First, compute a *normal step* toward minimizing constraint violation

$$v(x) = \frac{1}{2} \|c(x)\|^2 \Rightarrow \begin{cases} \min_{s^n \in \mathbb{R}^n} m_k^v(s^n) \\ \text{s.t. } \|s^n\| \leq \delta_k^v \end{cases}$$

- ▶ Second, compute multipliers y_k (or take from previous iteration).
- ▶ Third, compute a *tangential step* toward optimality:

$$\min_{s^t \in \mathbb{R}^n} m_k^f(s_k^n + s^t) \quad \text{s.t. } J_k s^t = 0, \quad \|s_k^n + s^t\| \leq \delta_k^f.$$

Main idea

Two-phase method combining trust funnel and TRACE.

- ▶ Trust funnel for globalization
- ▶ TRACE for good complexity bounds

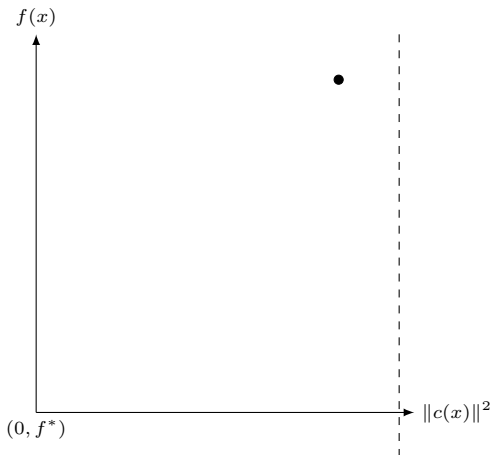
Phase 1 towards feasibility, two types of iterations:

- ▶ F-ITERATIONS improve objective and reduce constraint violation.
- ▶ V-ITERATIONS reduce constraint violation.

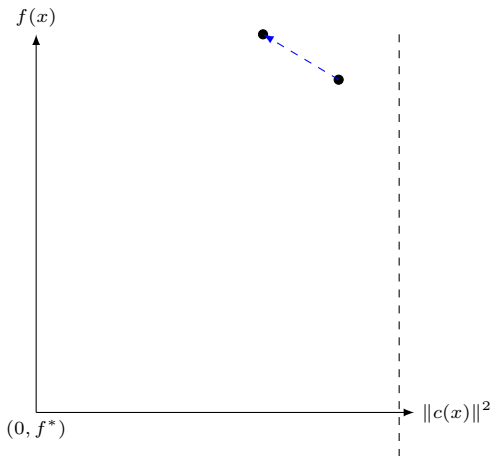
Our algorithm vs. basic trust funnel

- ▶ modified F-ITERATION conditions and a different funnel updating procedure
- ▶ uses TRACE ideas (for radius updates) instead of tradition trust region
- ▶ after getting approximately feasible, switches to “phase 2”.

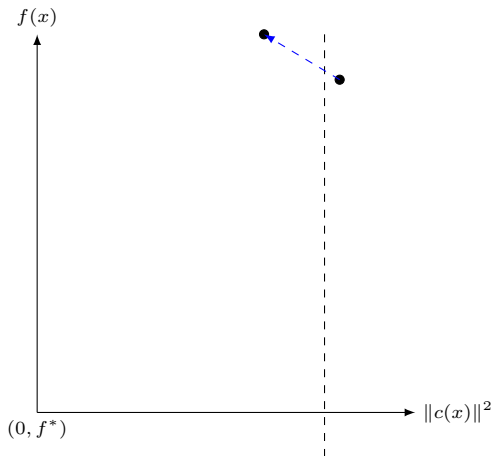
Our algorithm-Illustration



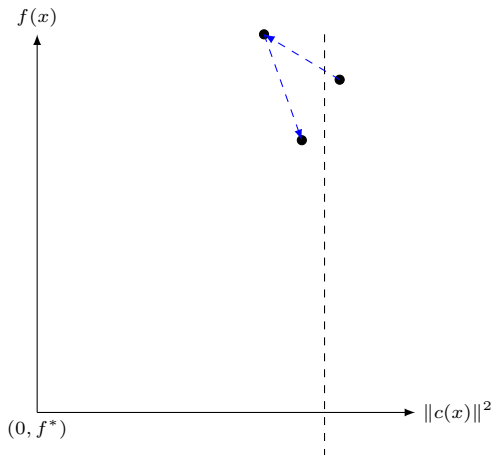
Our algorithm-Illustration



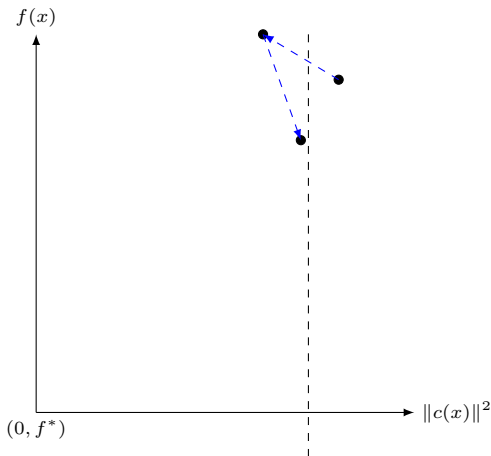
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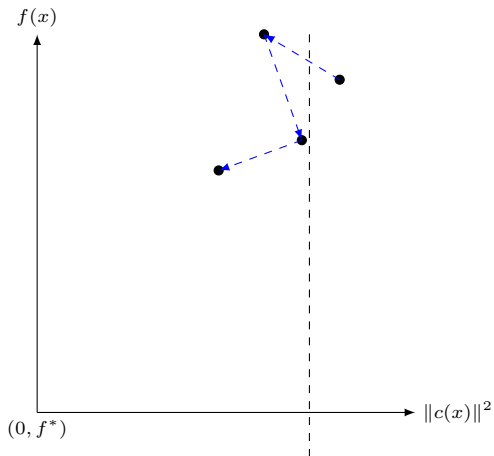
Our algorithm-Illustration



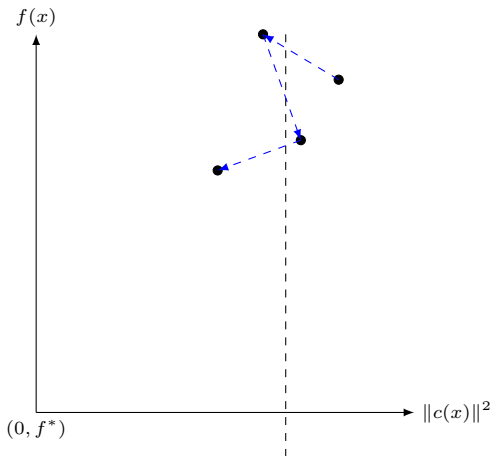
Our algorithm-Illustration



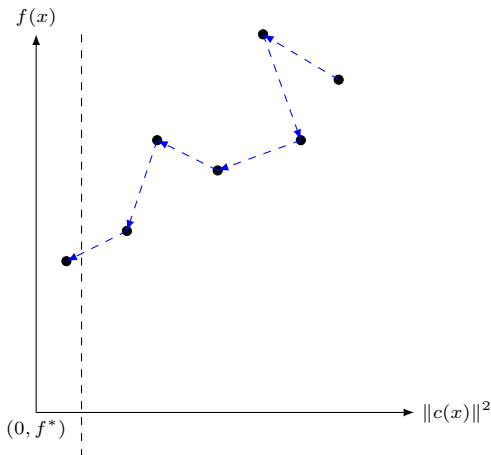
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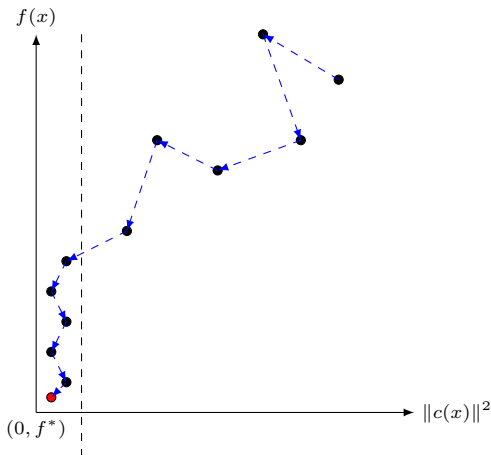
Our algorithm-Illustration



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Our algorithm-Illustration



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Phase 1

Recall that $\nabla v(x) = J(x)^T c(x)$ and define the iteration index set

$$\mathcal{I} := \{k \in \mathbb{N} : \|J_k^T c_k\| > \epsilon_v\}.$$

Theorem

For any $\epsilon_v \in (0, \infty)$, the cardinality of \mathcal{I} is at most $K(\epsilon_v) \in \mathcal{O}(\epsilon_v^{-3/2})$:

- ▶ $\mathcal{O}(\epsilon_v^{-3/2})$ successful steps and
- ▶ finite contraction and expansion steps between successful steps.

Corollary

If $\{J_k\}$ have full row rank with singular values bounded below by $\xi \in (0, \infty)$, then

$$\mathcal{I}_c := \{k \in \mathbb{N} : \|c_k\| > \epsilon_v/\xi\}$$

has cardinality $\mathcal{O}(\epsilon_v^{-3/2})$.

Phase 2

Options for phase 2:

- ▶ trust funnel method (no complexity guarantees) or
- ▶ “target-following” approach similar to Short-Step ARC to minimize

$$\Phi(x, t) = \|c(x)\|^2 + |f(x) - t|^2.$$

Theorem

For $\epsilon_f \in (0, \epsilon_v^{1/3}]$, the number of iterations until

$$\|g_k + J_k^T y\| \leq \epsilon_f \|(y_k, 1)\| \text{ or } \|J_k^T c_k\| \leq \epsilon_f \|c_k\|$$

is $\mathcal{O}(\epsilon_f^{-3/2} \epsilon_v^{-1/2})$.

Same complexity as Short-Step ARC:

- ▶ If $\epsilon_f = \epsilon_v^{2/3}$, then overall $\mathcal{O}(\epsilon_v^{-3/2})$
- ▶ If $\epsilon_f = \epsilon_v$, then overall $\mathcal{O}(\epsilon_v^{-2})$

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Implementation

MATLAB implementation:

- ▶ Phase 1: our algorithm vs. one doing V-ITERATION only
- ▶ Phase 2: trust funnel method [Curtis, Gould, Robinson, & Toint (2016)]

Termination conditions:

- ▶ Phase 1:

$$\|c_k\|_\infty \leq 10^{-6} \max\{\|c_0\|_\infty, 1\} \quad \text{or} \quad \begin{cases} \|J_k^T c_k\|_\infty \leq 10^{-6} \max\{\|J_0^T c_0\|_\infty, 1\} \\ \text{and } \|c_k\|_\infty > 10^{-3} \max\{\|c_0\|_\infty, 1\} \end{cases}$$

- ▶ Phase 2

$$\|g_k + J_k^T y_k\|_\infty \leq 10^{-6} \max\{\|g_0 + J_0^T y_0\|_\infty, 1\}.$$

Test set

Equality constrained problems (190) from CUTEst test set:

78	constant (or null) objective
60	time limit
13	feasible initial point
3	infeasible phase 1
2	function evaluation error
1	small stepsizes (less than 10^{-40})

Remaining set consists of 33 problems.

Problem	n	m	TF						TF-V-ONLY					
			Phase 1				Phase 2		Phase 1				Phase 2	
			#V	#F	f	$\ g + J^T y\ $	#V	#F	#V	f	$\ g + J^T y\ $	#V	#F	
BT1	2	1	4	0	-8.02e-01	+4.79e-01	0	139	4	-8.00e-01	+7.04e-01	7	136	
BT10	2	2	10	0	-1.00e+00	+5.39e-04	1	0	10	-1.00e+00	+6.74e-05	1	0	
BT11	5	3	6	1	+8.25e-01	+4.84e-03	2	0	1	+4.55e+04	+2.57e+04	16	36	
BT12	5	3	12	1	+6.19e+00	+1.18e-05	0	0	16	+3.34e+01	+4.15e+00	4	8	
BT2	3	1	22	8	+1.45e+03	+3.30e+02	3	12	21	+6.14e+04	+1.82e+04	0	40	
BT3	5	3	1	0	+4.09e+00	+6.43e+02	1	0	1	+1.01e+05	+8.89e+02	0	1	
BT4	3	2	1	0	-1.86e+01	+1.00e+01	20	12	1	-1.86e+01	+1.00e+01	20	12	
BT5	3	2	15	2	+9.62e+02	+2.80e+00	14	2	8	+9.62e+02	+3.83e-01	3	1	
BT6	5	2	11	45	+2.77e-01	+4.64e-02	1	0	14	+5.81e+02	+4.50e+02	5	59	
BT7	5	3	15	6	+1.31e+01	+5.57e+00	5	1	12	+1.81e+01	+1.02e+01	19	28	
BT8	5	2	50	26	+1.00e+00	+7.64e-04	1	1	10	+2.00e+00	+2.00e+00	1	97	
BT9	4	2	11	1	-1.00e+00	+8.56e-05	1	0	10	-9.69e-01	+2.26e-01	5	1	
BYRDSPHR	3	2	29	2	-4.68e+00	+1.28e-05	0	0	19	-5.00e-01	+1.00e+00	16	5	
CHAIN	800	401	9	0	+5.12e+00	+2.35e-04	3	20	9	+5.12e+00	+2.35e-04	3	20	
FLT	2	2	15	4	+2.68e+10	+3.28e+05	0	13	19	+2.68e+10	+3.28e+05	0	17	
GENHS28	10	8	1	0	+9.27e-01	+5.88e+01	0	0	1	+2.46e+03	+9.95e+01	0	1	
HS100LNP	7	2	16	2	+6.89e+02	+1.74e+01	4	1	5	+7.08e+02	+1.93e+01	14	3	
HS111LNP	10	3	9	1	-4.78e+01	+4.91e-06	2	0	10	-4.62e+01	+7.49e-01	10	1	
HS27	3	1	2	0	+8.77e+01	+2.03e+02	3	5	1	+2.54e+01	+1.41e+02	11	34	
HS39	4	2	11	1	-1.00e+00	+8.56e-05	1	0	10	-9.69e-01	+2.26e-01	5	1	
HS40	4	3	4	0	-2.50e-01	+1.95e-06	0	0	3	-2.49e-01	+3.35e-02	2	1	
HS42	4	2	4	1	+1.39e+01	+3.94e-04	1	0	1	+1.50e+01	+2.00e+00	3	1	
HS52	5	3	1	0	+5.33e+00	+1.54e+02	1	0	1	+8.07e+03	+4.09e+02	0	1	
HS6	2	1	1	0	+4.84e+00	+1.56e+00	32	136	1	+4.84e+00	+1.56e+00	32	136	
HS7	2	1	7	1	-2.35e-01	+1.18e+00	7	2	8	+3.79e-01	+1.07e+00	5	2	
HS77	5	2	13	30	+2.42e-01	+1.26e-02	0	0	17	+5.52e+02	+4.54e+02	3	11	
HS78	5	3	6	0	-2.92e+00	+3.65e-04	1	0	10	-1.79e+00	+1.77e+00	2	30	
HS79	5	3	13	21	+7.88e-02	+5.51e-02	0	2	10	+9.70e+01	+1.21e+02	0	24	
MARATOS	2	1	4	0	-1.00e+00	+8.59e-05	1	0	3	-9.96e-01	+9.02e-02	2	1	
MSS3	2070	1981	12	0	-4.99e+01	+2.51e-01	50	0	12	-4.99e+01	+2.51e-01	50	0	
MWRIGHT	5	3	17	6	+2.31e+01	+5.78e-05	1	0	7	+5.07e+01	+1.04e+01	12	20	
ORTHREGB	27	6	10	15	+7.02e-05	+4.23e-04	0	6	10	+2.73e+00	+1.60e+00	0	10	
SPIN20P	102	100	57	18	+2.04e-08	+2.74e-04	0	1	time	+1.67e+01	+3.03e-01	time	time	

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BT11	5	3	6	1	+8.25e-01	+4.84e-03	2	0	1	+4.55e+04	+2.57e+04	16	36	
BT12	5	3	12	1	+6.19e+00	+1.18e-05	0	0	16	+3.34e+01	+4.15e+00	4	8	

Summary of results

Our algorithm, at the end of phase 1

- ▶ for 26 problems, reaches a smaller function value
- ▶ for 6 problems, reaches the same function value

Total number of iterations of our algorithm

- ▶ for 18 problems is smaller
- ▶ for 8 problems is equal

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- ▶ Proposed an algorithm for equality constrained optimization
- ▶ Trust funnel algorithm with improved complexity properties
- ▶ Promising performance in practice based on our preliminary numerical experiment
- ▶ A step toward practical algorithms with good iteration complexity

F. E. Curtis, D. P. Robinson, and M. Samadi. Complexity Analysis of a Trust Funnel Algorithm for Equality Constrained Optimization. Technical Report 16T-013, COR@L Laboratory, Department of ISE, Lehigh University, 2016.